

LangSec for AppSec folks



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Day 2025

Dr. Lars Hermerschmidt



Product Owner Security Engineering



Industry

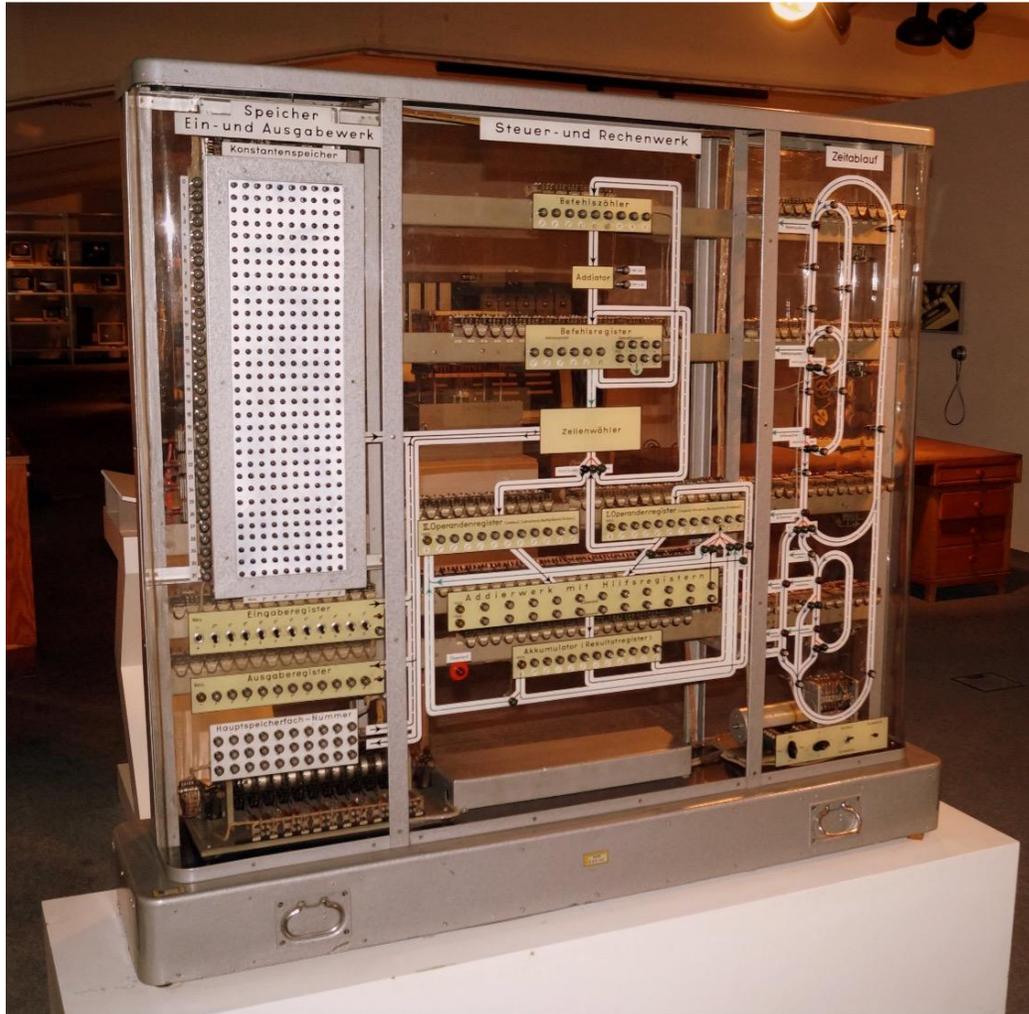
- Building Security Champions programs since 2017
- Coach and trainer for software security
- Help in DevOps transformation



Research

- LangSec, eliminating injections
- Threat Modeling automation
- Evidence based software security programs
 - security capability model Security Belts





Software is Input

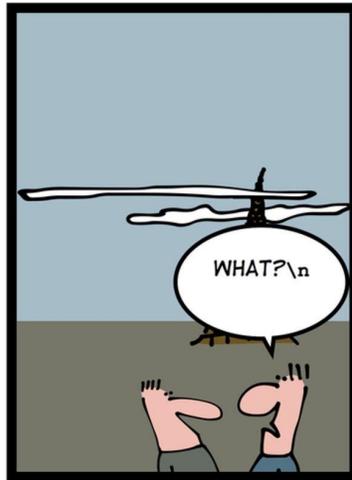
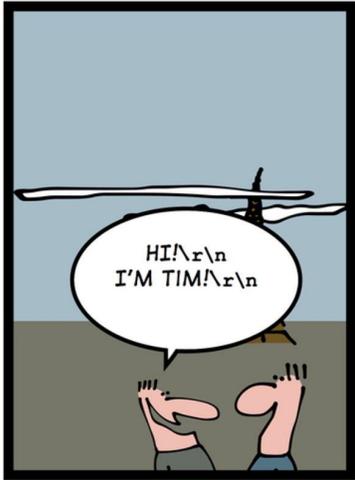
Von Neumann architecture

- General purpose machine
- Software is loaded and executed by the machine
- Software is stored like data

Input drives Software

- „Datenverarbeitung“ is an illusion
- Software reacts on input
- Attackers can use all offered processing capabilities

The View from the Tower of Babel



BABYLON

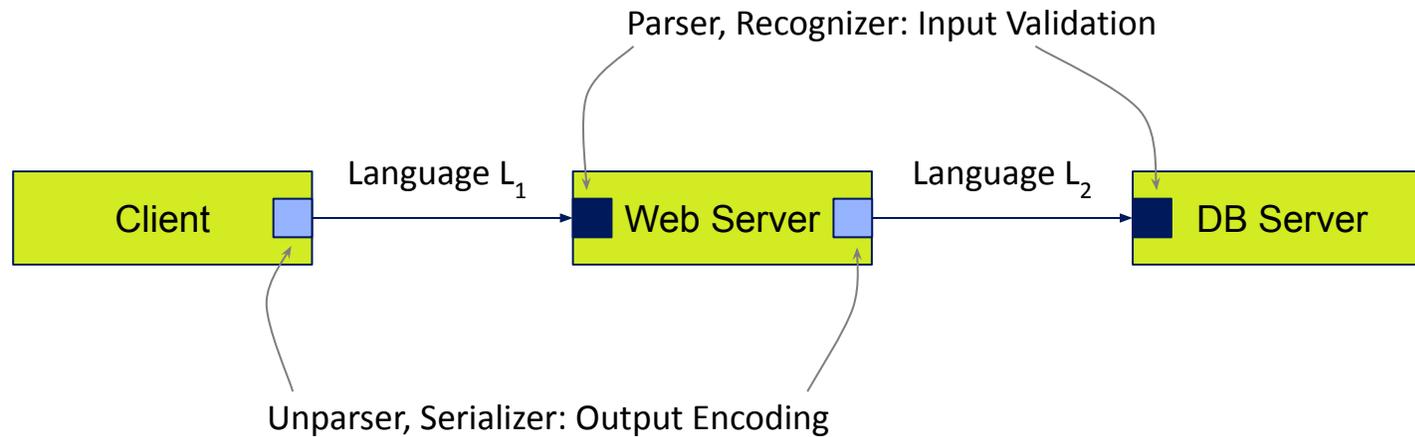
LangSec attributes the insecurity of software to ad hoc programming of input handling

- Treat all input as formal language
- Input handling code is a (un)parser for that language
- That (un)parser implementation must match the grammar used to define the language

Constructing secure (un)parsers is as complex as secure cryptography, so reuse known good libraries



Definitions - Where is that Language?





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01 Parser, Recognizer: Input Validation



Common Bad Advices: “Don’t trust input data”

Input Validation

- Validate all inputs against a specification
- Accept only those input fulfilling the specification
- Reject all others, NEVER try to guess/heal bad input

Input specification

Reject Known bad with deny list or black list: $^{\wedge} [^{\wedge} \mathbf{a-z}] + \$$

Only accept known inputs with whitelist: $^{\wedge} [0-9] + \$$

Beware: Regular Expressions can only recognize regular languages (no nesting)



FX: Blitzableiter (2010)

Security Concerns with Adobe Flash

Adobe Flash Attack Surface

- Flash files (SWF) is a container format for:
 - Vector graphics data (shapes, morphing, gradients)
 - Pixel graphics formats (various JPEG, lossless bitmaps)
 - Fonts and text
 - Sound data (ADPCM, MP3, Nellymoser, Speex)
 - Video data (H.263, Screen Video, Screen Video V2, On2 Truemotion VP6)
 - Virtual machine byte code for the Adobe Virtual Machines (AVM)
- All data structures from file format version 3 until the current version 10 are still supported
- The parser is completely written in unmanaged languages (C/C++)

**FX exploited countless vulnerabilities in Flash
Adobe did not manage to patch them**

**Blitzableiter a rigor parser used in front of the
vulnerable Adobe Flash Plugin:**

1. Recognize only valid Flash
2. Create a new Flash file from the accepted content
3. Pass the new Flash file to the Flash Plugin

Blitzableiter defended all exploitation attempts

FX found LangSec

**The technique is now called Content Disarm and
Reconstruction**



Antipattern: Shotgun Parser

```
186 /**
187  * This will parse the stream and populate the PDDocument object. This will close the
188  * done parsing.
189  *
190  * @param lenient activate leniency if set to true
191  * @throws InvalidPasswordException If the password is incorrect.
192  * @throws IOException If there is an error reading from the stream or corrupt data is
193  */
194 public PDDocument parse(boolean lenient) throws IOException
195 {
196     setLenient(lenient);
197     // set to false if all is processed
198     boolean exceptionOccurred = true;
199     try
200     {
201         // PDFBOX-1922 read the version header and rewind
202         if (!parsePDFHeader() && !parseFDFHeader())
203         {
204             throw new IOException( "Error: Header doesn't contain versioninfo" );
205         }
206
207         if (!initialParseDone)
208         {
209             initialParse();
210         }
211         exceptionOccurred = false;
212         PDDocument pdDocument = createDocument();
213         pdDocument.setEncryptionDictionary(getEncryption());
214         return pdDocument;
215     }
216     finally
217     {
218         if (exceptionOccurred && document != null)
219         {
220             IOUtils.closeQuietly(document);
221             document = null;
222         }
223     }
224 }
```

- Hand-written parser
- The recognition logic is spread all over the application
- State is allocated before the input is fully recognized



- Apache Tika uses Apache PDFBox Parser
- Hand-written parser
- Does it comply to some PDF grammar?

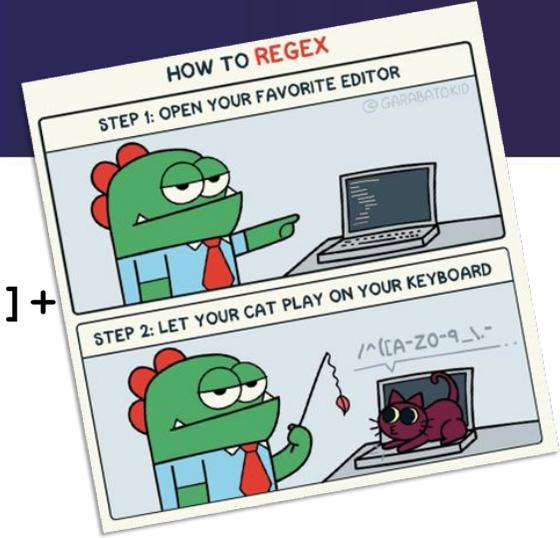
<https://svn.apache.org/viewvc/pdfbox/trunk/pdfbox/src/main/java/org/apache/pdfbox/pdfparser/PDFParser.java?revision=1878544&view=co>



**Don't implement parsers in
\$YourFavoredProgrammingLanguage!**



Writing Secure Parsers is hard



`^([a-zA-Z0-9])(([\-.]|[_]+)?([a-zA-Z0-9]+))*(@){1}[a-z0-9]+,3})|([a-z]{2,3}[.]{1}[a-z]{2,3}))$`

Regular Expression Catastrophic Backtracking Denial of Service

Polyglot: A files that is accepted by multiple parsers as different files

Parser Differential Bug

- Check with parser A
- Execute with parser B

github.com/Polydet/polyglot-database

README MIT license

Polyglots database

This repository is a data set of polyglot files. It's purpose is to be used to test
Please note that this repository uses <https://git-lfs.github.com/>.

	7z	afsk	agc	apk	avi	bmp
7ZIP+JAR.7z	x					
7ZIP+RAR-1.7z	x					
7ZIP+RAR-2.7z	x					
7ZIP+ZIP.7z	x					
AVI+HTML.avi					x	
AVI+ZIP.avi					x	
BMP+HTML+JAR.bmp						x
BMP+JAR.bmp						x
DOCX+ELF+JAR+PDF+RAR-1.zip						

github.com/corkami/mitra

README MIT license

Mitra

A tool to generate weird files (parasites, polymocks, polyglots, near-polyglots, crypto-polyglots) for many formats.
Loosely named after [Μιθραδάτης](#), a polyglot historical figure (it's pronounced `mitra`).

[What's new.](#)

How to use

- `mitra.py file1.pdf /dev/null -f --pad 2` gives you a standard PDF with an extra range of arbitrary data.
- `mitra.py file1.png file2.dcm` gives you a working PNG/DICOM polyglot.

Check Corkami [mini](#) or [tiny](#) PoCs for input files, and the formats [repository](#) for some extra technical info.

Features

It tries different layouts: Stacks (appended data), Cavities (blank space), Parasites (comments), Zippers (mutual comments).

It stores the offsets where the payloads 'switch sizes' between parenthesis for ambiguous ciphertexts.



Standardized language

- Find existing good (Un)Parser
- Review Parser with LangSec knowledge
- Work on Parsetree to access the data

Your own language

- Define language through grammar
- Generate (Un)Parser or use (Un)Parser Combinator
- Work on Parsetree to access the data



Don't fear the grammar

CSV Example Grammar

```
29  grammar CSV;
30
S 31  csvFile: hdr row+ ;
32  hdr: row ;
33
P 34  row: field (',' field)* '\r'? '\n' ;
35
36  field
37  : TEXT
38  | STRING
39  |
40  ;
41
42  TEXT  : ~[, \n\r"]+ ;
43  STRING: '"' (""|~'')* '"' ; // quote-quote is an escaped quote
```

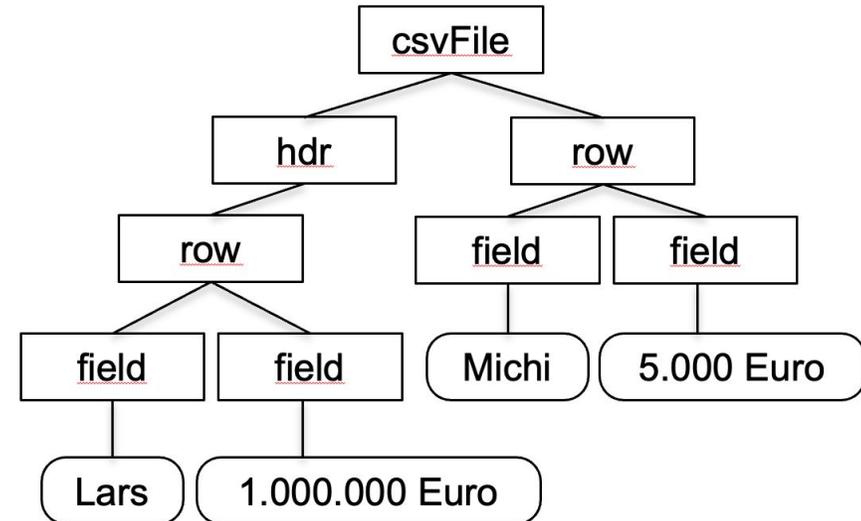
Grammar $G = (N, T, P, S)$
Nonterminal Symbols
Terminal Symbols
Productions
Start Symbol

<https://github.com/antlr/grammars-v4/blob/master/csv/CSV.g4>



CSV Example

Lars,1.000.000 Euro
Michi,5.000 Euro





Chomsky Hierarchy for Transport Languages

allgemein, Typ 0, Produktionen: beliebig, Automat: Turing-Maschine

kontextsensitiv, Typ 1, Produktionen: $\alpha \rightarrow \beta$, $|\alpha| \leq |\beta|$, Automat: linear beschränkte TM

kontextfrei, Typ 2, Produktionen: $A \rightarrow \alpha$, Automat: Kellerautomat

regulär, Typ 3, Produktionen: $A \rightarrow a \mid aB$, Automat: endlicher Automat

Context free Grammar: Halting problem decidable and computable

- „One can construct a parser and proof it's correctness.“
- Proof is challenging (ASN.1 compiler has been verified using Coq)

Above Context free Grammar = Remote Code Execution

- Parser has to be a Turing Machine
- Natural Languages are mildly context-sensitive = LLMs are doomed



Computational Cliff

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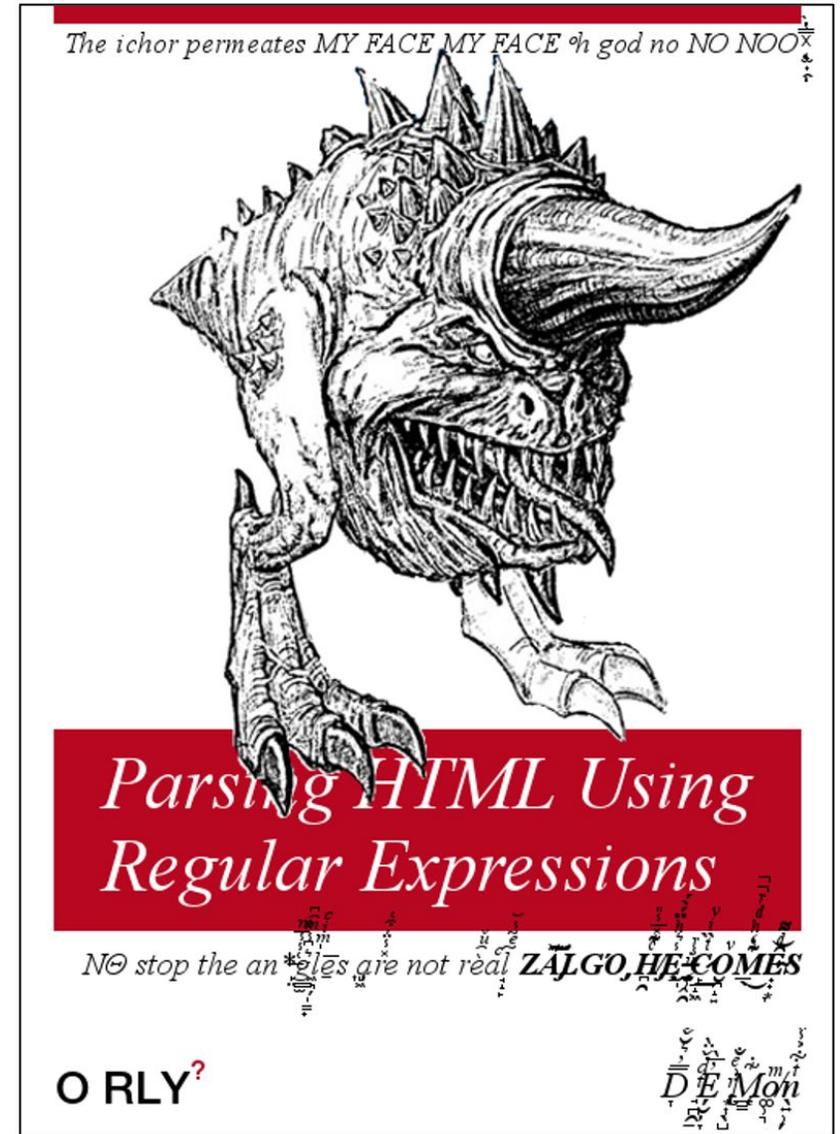
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Never use regular Expressions to parse more complex languages





Do not specify formats that are more complex than context free



Example: XML External Entity Injection (XXE)

- XML parsers load external entities from filesystem or network 🙄 (language designers thought that would be a good idea)
- A Webservice processing this XML would return the /etc/passwd file to the caller

```
1  <?xml version="1.0" encoding="ISO-8859-1"?>
2  <!DOCTYPE foo [
3    <!ELEMENT foo ANY >
4    <!ENTITY xxe SYSTEM "file:///etc/passwd" >]>
5  <foo>&xxe;</foo>
```

Syntax

- Parsing enforces syntactical correctness on input
- Programs can decide which part of that input to process
 - Filtering of instructions is possible

Semantics

- Processing Input results in semantic action
 - Execute parsed instructions on Turing Machine
- (unintentionally) exposing semantic actions may lead to vulnerabilities, e.g., Log4Shell



Rice's theorem, Code Scanners

Programming Languages are awesome

- Turing complete
- Open to solve all problems

Programming Languages are scary

- Turing complete
- Unintended functionality possible

Rice's theorem: “non-trivial, semantic properties of programs are undecidable”

It is impossible to predict what code or scripts are doing before executing them.

Implications for programming languages

- All static code analysis searching for bugs have to solve an undecidable problem and will fail.
- Dynamic analysis, i.e., security scanners is no better.



Rice's theorem, Code Scanners, and Antivirus

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Implications for file formats

- Antivirus does not work.
- Executing in a sandbox and observing is no better.
- Solution: Remove semantic actions, i.e., scripts from file formats.



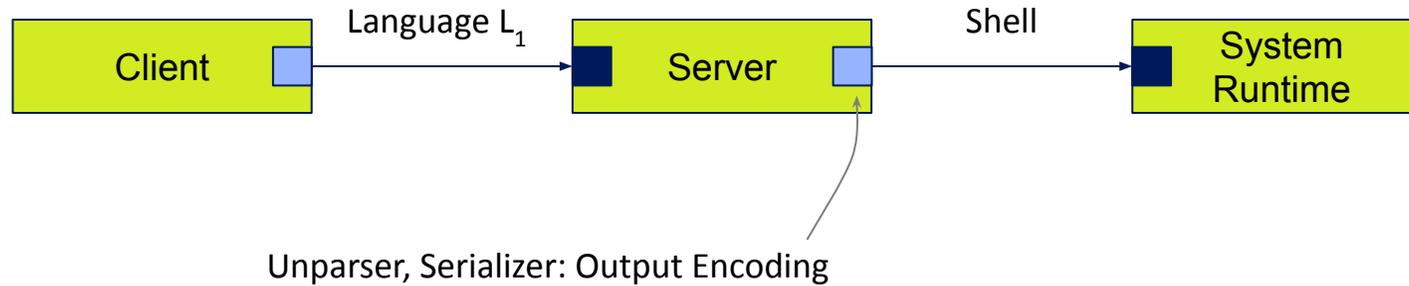
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02 Unparser, Serializer: Output Encoding





Protection Against Input during Output Creation

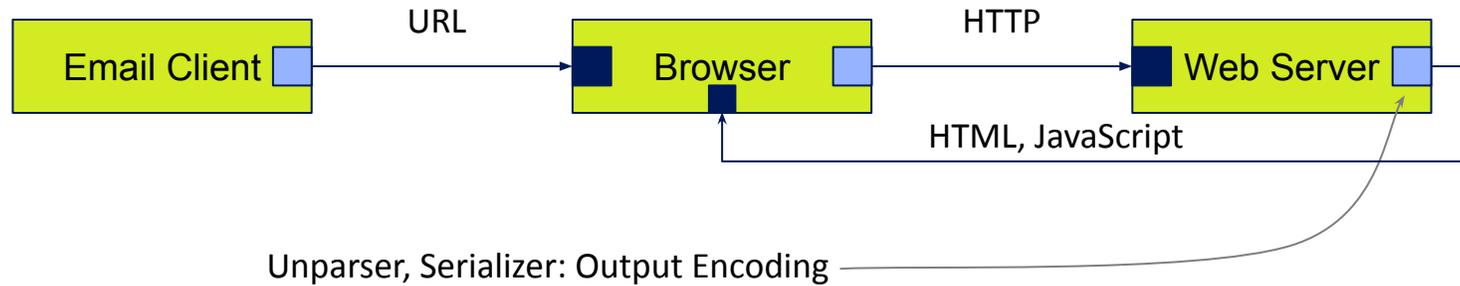


```
1 String btype = request.getParameter("backuptype");
2 String cmd = new String("cmd.exe /K \"c:\\util\\rmanDB.bat \"
3 +btype+
4 "&&c:\\utl\\cleanup.bat\"")
5
6 System.Runtime.getRuntime().exec(cmd);
```

<https://cwe.mitre.org/data/definitions/77.html>



Protection Against Input during Output Creation

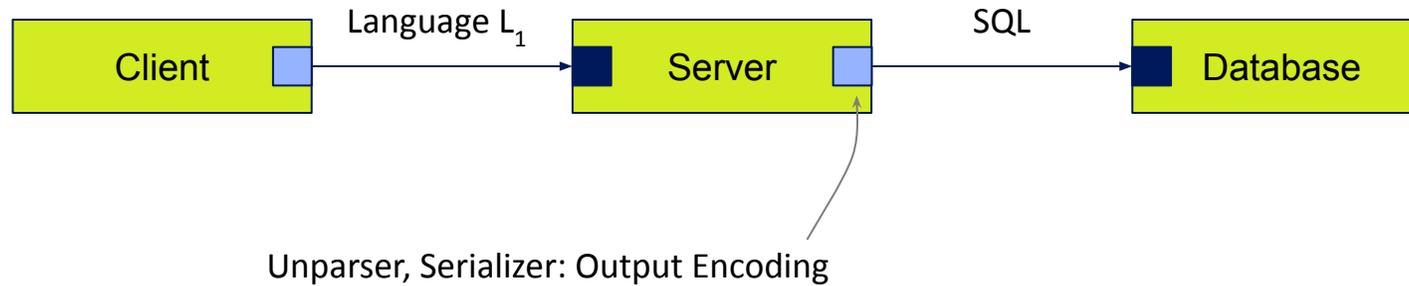


```
1 $username = $_GET['username'];  
2 echo '<div class="header"> Welcome, ' . $username . '</div>';
```

<https://cwe.mitre.org/data/definitions/79.html>



Protection Against Input during Output Creation



```
1 String firstname = req.getParameter("firstname");
2 String lastname = req.getParameter("lastname");
3
4 String query = "SELECT id, firstname, lastname FROM authors WHERE "
5 | + "firstname = "+firstname+" and lastname = "+lastname;
6 PreparedStatement pstmt = connection.prepareStatement( query );
7 try
8 {
9 |   ResultSet results = pstmt.execute( );
10 }
```

https://owasp.org/www-community/attacks/SQL_injection

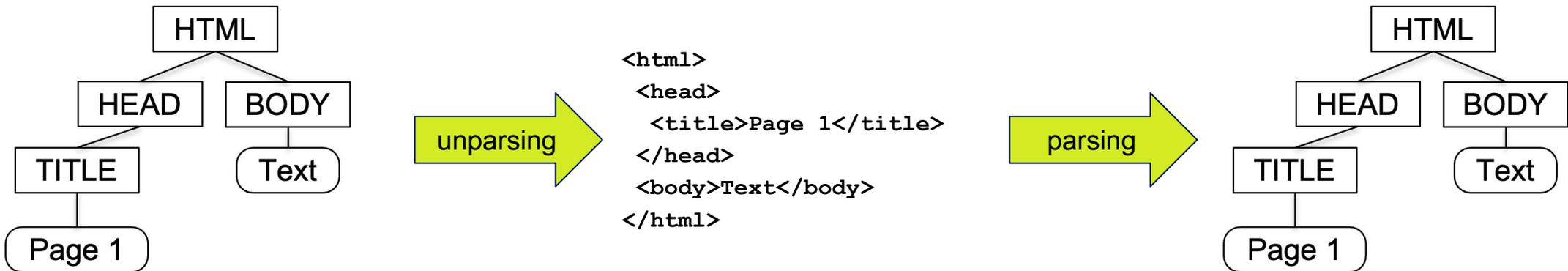


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String concatenation leads to Injections



(Un)Parser Roundtrip





(Un)Parser Roundtrip





Injections from a Language Perspective

“malicious input” is made of keywords from the grammar

Keywords separate data from code

Injection vulnerabilities are rooted in language design

All non-trivial languages with keywords are injectable

Languages using length field seem to be secure

- “Total Length” in IPv4 Header
- IEEE 802.3 MAC

```
29 grammar CSV;
30
31 csvFile: hdr row+ ;
32 hdr : row ;
33
34 row : field (',' field)* '\r'? '\n' ;
35
36 field
37     : TEXT
38     | STRING
39     ;
40
41
42 TEXT  : ~[, \n\r"]+ ;
43 STRING : '"' ( '""' | ~'"')* '"' ; // quote-quote is an escaped quote
```

Keywords (part of T)

<https://github.com/antlr/grammars-v4/blob/master/csv/CSV.g4>



Vulnerable CSV Implementation

```
8  @RestController
9  public class BadController {
10
11     @GetMapping("/bad/contributions/{name}")
12     public String getContribution(@PathVariable("name") String name) {
13         return name + ',' + "10.000 Euro";
14     }
15
16     @GetMapping("/bad/bill")
17     public String printBill(@RequestParam("contribution") String contribution) {
18         String[] fields = contribution.split(regex: ",");
19
20         StringBuilder out = new StringBuilder();
21         out.append(str: "----- eBon -----\n");
22         out.append("Name: " + fields[0] + "\n");
23         out.append("Sum: " + fields[1] + "\n");
24         return out.toString();
25     }
26
27 }
```

Insecure unparser: String concatenation

Bad parser: Regular Expression for a context free language



Good CSV Implementation

```
14 @RestController
15 public class GoodController {
16
17     @GetMapping("/good/contributions/{name}")
18     public String getContribution(@PathVariable("name") String name) throws IOException {
19         StringBuilder out = new StringBuilder();
20         try (CSVPrinter printer = new CSVPrinter(out, CSVFormat.DEFAULT)) {
21             printer.printRecord(name, "10.000 Euro");
22         }
23         return out.toString();
24     }
25
26     @GetMapping("/good/bill")
27     public String printBill(@RequestParam("contribution") String contribution) throws IOException {
28         CSVParser reader = CSVFormat.DEFAULT.parse(new StringReader(contribution));
29
30         StringBuilder out = new StringBuilder();
31         reader.forEach(record -> {
32             out.append(str: "----- eBon -----\n");
33             out.append("Name: " + record.get(0) + "\n");
34             out.append("Sum: " + record.get(1) + "\n");
35         });
36         return out.toString();
37     }
38 }
```

Unparser from a library
Work with the Unparser API

Parser from a library
Get data from the parser API



Recipe for Secure Input Handling

Use (Un)Parser Libraries that do encoding automatically correct

- Don't write encoding code yourself
- Context-sensitive Encoding, e.g. in HTML, is too complex to get it right

When using an existing language (HTML, SQL, CSV, ...)

- Find existing good (Un)Parser.
 - Fuzz-Test encoding
 - Review (Un)Parser with LangSec knowledge (Shotgun vs. Grammar based)
- Work on Parsetree to access the data

When creating your own communication language: Really? Why not JSON, or Cap'n Proto?

- Define language through grammar
- Generate (Un)Parser
- Work on Parsetree to access the data



Solving LangSec in Programming Languages

Interact with a language

- Make using good (un)parsers for existing languages easy
- Make constructing (un)parsers easy
- Build in (un)parser construction capabilities into programming languages core

Remove byte operations

- Remove simple Input Output functions like byte wise read and println()
- Remove string concatenation



Meanwhile at Google

SafeCoding - Eliminate Bug Classes

Bug Classes

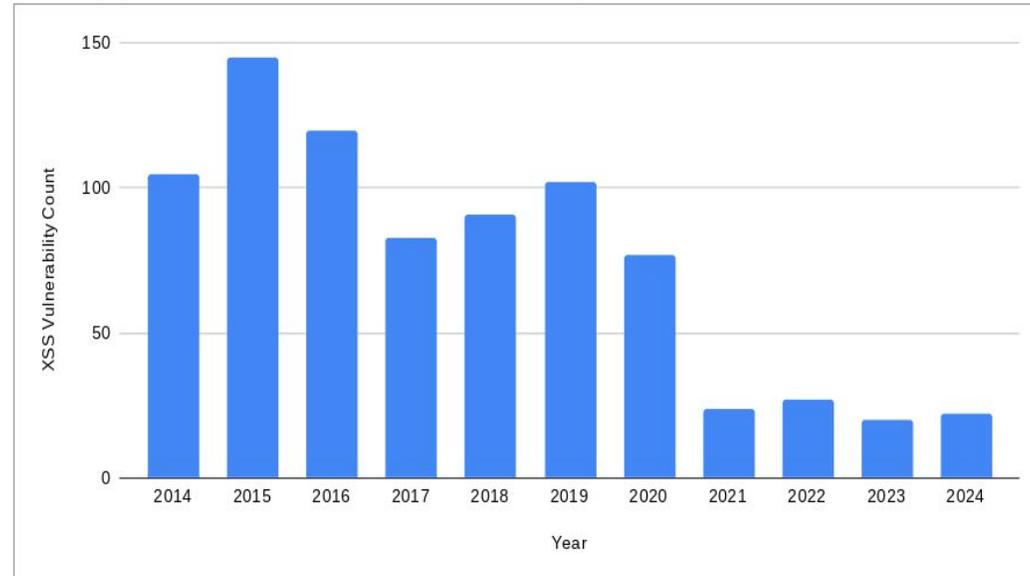
Encryption: Easy to build an application without encryption-in-transit

Injection: The web's core building blocks (HTML, JS, URLs) allow easily mixing code and data

Isolation: Cross-origin interactions between different applications are allowed by default

Solution

Providing a platform to developers that is Safe





AppSec should learn from LangSec

Input Validation is a misleading term

- Correct unparsing prevents Injections
- Strict parsing does not solve the root cause of injections, it makes programs robust against input

Computers (including AI) can not determine what programs do

- SAST and DAST try to solve an undecidable problem and fail.
- Antivirus does not work.

Eliminate Bug Classes

Communication language above deterministic context free introduce Remote Code Execution by design

“Rolling your own (un)parser is like rolling your own crypto algo”